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Paul Baran, Network Theory, and the Past, Present, and Future of the Internet

Christopher S. Yoo*

Paul Baran’s seminal 1964 article “On Distributed Communications Networks” that first proposed packet switching also advanced an underappreciated vision of network architecture: a lattice-like, distributed network, in which each node of the Internet would be homogeneous and equal in status to all other nodes. Scholars who have subsequently embraced the concept of a lattice-like network approach have largely overlooked the extent to which it is both inconsistent with network theory (associated with the work of Duncan Watts and Albert-László Barabási), which emphasizes the importance of short cuts and hubs in enabling networks to scale, and the actual way, the Internet initially deployed, which relied on a three-tiered, hierarchical architecture that was actually what Baran called a decentralized network. However, empirical studies reveal that the Internet’s architecture is changing: it is in the process of becoming flatter and less hierarchical, as large content providers build extensive wide area networks and undersea cables to connect directly to last-mile networks. This change is making the network more centralized rather than becoming more distributed. As a result, this article suggests that the standard reference model that places backbones at the center of the architecture should be replaced with a radically different vision: a stack of centralized star networks, each centered on one of the leading content providers.

INTRODUCTION

Paul Baran’s article “On Distributed Communications Networks”1 represents an underappreciated landmark in the history of the Internet.2 It articulated a bold new approach to

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2 A Westlaw search reveals only fifteen citations to Baran’s article and research memoranda, mostly consisting of passing references or applications to other areas of the law. I am aware of only two law review articles that reproduce Baran’s classic figure; of these, one applies the principles to securities law, and the other deletes the
communications networking that would serve as one of the foundations for the Internet. Baran’s proposal was complex, but its essence can be distilled into three basic principles. First, to make the system robust against physical attack, all nodes would be connected to the network by multiple links. Second, communications would be divided into individual packets that would travel independently through the network and be reassembled once they reached their destinations. Third, to take advantage of multiple links and to eliminate the vulnerability associated with a single centralized point of control, each node would route traffic independently on a simple store-and-forward basis.

Though controversial when initially proposed, the second and third principles ultimately laid the basis for the packet switching method that would eventually become one of the Internet’s foundational technologies. The first principle, on the other hand, proved to be less influential: it was neither incorporated into the design of the Advanced Research Projects Agency Network (ARPANET) nor the subsequent Internet.

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decentralized network that ends up being the most important of the three network architectures. See Patrick S. Ryan, Wireless Communications and Computing at a Crossroads: New Paradigms and Their Impact on Theories Governing the Public’s Right to Spectrum Access, 3 J. ON TELECOMM. & HIGH TECH. L. 239, 262 fig.1 (2005); Will Schildknecht, Comment, Designing for Robustness: Overcoming Systemic Risk in the Political Branches, 103 CALIF. L. REV. 433, 440 fig.1 (2015). Baran’s diagram is occasionally reproduced in histories of the Internet. E.g., KATIE HAFNER & MATTHEW LYON, WHERE WIZARDS STAY UP LATE 58 (1996). It also appears in leading books on network science. E.g., ALBERT-LÁSZLÓ BARABÁSI, LINKED: HOW EVERYTHING IS CONNECTED TO EVERYTHING ELSE AND WHAT IT MEANS FOR BUSINESS, SCIENCE, AND EVERYDAY LIFE 145 fig.11.1 (2002); MARK BUCHANAN, NEXUS: SMALL WORLDS AND THE GROUNDBREAKING THEORY OF NETWORKS 79 fig.7 (2002).

3 Baran RAND Paper, supra note 1, at 2.
4 Id. at 23–27.
5 Id. at 27–30.
7 Id. at 44, 195 39–40 (noting that elements of Baran’s system “that were specifically adapted to the Cold War Threat, including very high levels of redundancy” were not adopted by Davies or by the ARPANET, while other features, such as “adaptive routing, and efficient packet switching,” were); BARABÁSI, supra note 2, at 144, 147 (noting that “the topology of today’s Internet has little to do with [Baran’s distributed] vision” and that “[i]ronically, the principles directing today’s Internet match Baran’s original vision in every respect except the guiding principle that motivated his work: undercutting vulnerability to attacks”).
The misconception that Internet represents a lattice-like distributed network continues to resonate in many quarters. It is sometimes erroneously presented as the architecture of the current Internet. More recently, the supposed historical provenance of the distributed architecture has sometimes been invoked as validation for wireless mesh networks and other technologies that employ similar topologies. Even more problematically, emerging research investigating whether the Internet’s architecture is flattening often mischaracterizes the networks as becoming more mesh-like.

This Article attempts to give Baran’s first principle its proper due by acknowledging both its contributions and its limitations. First, this Article rebuts claims that the Internet, at its inception, incorporated a distributed topology by placing those arguments in historical context. Second, this Article draws on burgeoning research in network science to explain the critical role that hierarchical network structures play in ensuring that the network can operate in a scalable manner. Third, this Article reviews recent developments in the Internet’s architecture to demonstrate how the network is flattening. A careful analysis of the evidence about average path length reveals that the network is becoming increasingly centralized rather than more distributed.

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8 See BARRIE SOSINSKY, NETWORKING BIBLE 456 (2009) (“The Internet was designed to be a highly redundant mesh structure that could survive any outages to a substantial portion of the network and still be operable.”); Ryan, supra note 2, at 264 (“On the right side of Figure 1, the ‘distributed’ graphic depicts the way in which the Internet operates today.”).

9 See Yochai Benkler, Some Economics of Wireless Communications, 16 HARV. J.L. & TECH. 25, 75 (2002) (“The basic characteristic of the Internet’s robustness—its redundancy and decentralized architecture—is replicated in open wireless networks at the physical layer of the communications infrastructure”).


11 This analysis extends earlier research published in the predecessor of this journal. Christopher S. Yoo, Innovations in the Internet’s Architecture that Challenge the Status Quo, 8 J. ON TELECOMM. & HIGH TECH. L. 79 (2010).
I. Baran’s Unrealized Vision of a Distributed Networks

This Part describes Baran’s innovative reconceptualization of how to design a communications network and analyzes the extent to which his ideas were received. Section A lays out the motivation driving Baran’s vision and the principles underlying it. Section B explains why those ideas were not incorporated into the actual Internet.

A. Baran’s Vision of a Distributed Network

Baran’s motivation for creating his revolutionary approach to communications networking was driven by the dominant issue of the 1960s: the Cold War and the overhanging threat of nuclear annihilation. America’s nuclear capability depended on maintaining what the strategic defense community called “minimum essential communication,” which was the amount of connectivity needed for the U.S. to credibly maintain the threat of mutually assured destruction, which most observers believed was essential to deterring the Soviet Union from attempting a preemptive strike. Unfortunately, the existing communications infrastructure in the U.S. was dangerously fragile. Military leaders and analysts were concerned that an attack could cripple the country’s communications capabilities to the point where the President could not order a counterstrike.

12 Hafner & Lyon, supra note 2, at 55, 56 (noting that Baran felt “the problem of building a more stable communications infrastructure . . . was the most important work he could be doing” and quoting Baran as saying that his work “was done in response to the most dangerous situation that ever existed”).
13 Id. at 55.
14 Hafner & Lyon, supra note 2, at 55.
Accordingly, Baran focused his efforts on improving the survivability of the communications network.\textsuperscript{15} His principal mechanism for doing so was redundancy.\textsuperscript{16} The choice, as Baran framed it, came down to three possible network architectures: a \textit{centralized network}, a \textit{decentralized network}, and a \textit{distributed network}, as depicted in Figure 1.\textsuperscript{17} The centralized network is a classic star network, in which the peripheral nodes connect directly to a central node.\textsuperscript{18} The fact that destruction of the central node would disconnect every node in the system leaves this network extremely vulnerable to attack.\textsuperscript{19} The decentralized network is comprised of a number of smaller stars connected to form a larger star, with an added link connecting the two stars at the top of the graph to form a loop.\textsuperscript{20} This network architecture represents the then-current structure of the U.S. telephone network.\textsuperscript{21} Although more resilient than a centralized network to attack, a decentralized network can still be disabled by an attack on a small number of nodes.\textsuperscript{22} The distributed network, on the other hand, is more resilient. It is a lattice-like architecture in which the nodes are undifferentiated—they are all functionally similar and have a similar number of links to other nodes. Or, to use the language that will emerge later in this Article, the nodes have a similar enough number of connections to other nodes so that the network has a characteristic scale.\textsuperscript{23}

\textsuperscript{15} Abbate, \textit{supra} note 67, at 10–11; Hafner \& Lyon, \textit{supra} note 2, at 54–57.
\textsuperscript{16} See Figure 1; interestingly, some discussions omit the intermediate decentralized example. Ryan, \textit{supra} note 2, at 262 fig.1.
\textsuperscript{17} See graph (a) on the left-hand side of Figure 1.
\textsuperscript{18} Baran RAND Paper, \textit{supra} note 1, at 3.
\textsuperscript{19} See graph (b) on the left-hand side of Figure 1.
\textsuperscript{21} Baran RAND Paper, \textit{supra} note 1, at 3.
\textsuperscript{22} Barabási, \textit{supra} note 2, at 70.
For Baran, distributed networks rely on redundancy to increase the likelihood that two particular nodes could maintain contact with one another should other nodes in the network fail.24 Baran’s model indicates that connecting each node to only three or four others would be sufficient for a distributed network to approach the maximum theoretical level of survivability.25

Baran’s other principles followed as corollaries from his basic commitment to redundancy. For example, the fact that communications in a distributed network has to traverse an arbitrary number of hops meant that the network must necessarily be digital, since analog

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25 Id. at 3.
signals degrade every time they pass from one link to another.\textsuperscript{26} Furthermore, redundant links would not make the network more resilient unless each of the nodes were able to make independent routing decisions and reroute traffic autonomously in the face of partial network failure.\textsuperscript{27} That his design depended on nodes having sufficient intelligence to make their own routing decision meant that each would have to consist of a low-cost, unmanned computer.\textsuperscript{28}

Baran’s was not the first design to use a distributed architecture to increase survivability, but it was the first to use both unmanned, decentralized switching and an all-digital network with substantial intelligence concentrated in the nodes.\textsuperscript{29} Rather than attempting to maximize the reliability of each network element, Baran’s design relied on redundancy to compensate for any failures. Specifically, it used a large number of cheap, autonomous nodes to reinforce network survivability.\textsuperscript{30}

Baran’s design represented a radical departure from the traditional approach to networking, which relied on centralized control and hardening of key locations to ensure reliability and survivability.\textsuperscript{31} In the words of RAND’s history commemorating its work with the Air Force, Baran’s network was “unusual” in that each node had “equal status” and was capable of receiving, routing, and transmitting information autonomously.\textsuperscript{32}

\textsuperscript{26} ABBATE, \textit{supra} note 6, at 16; HAFNER & LYON, \textit{supra} note 2, at 57.
\textsuperscript{27} ABBATE, \textit{supra} note 6, at 13.
\textsuperscript{28} \textit{Id.} at 16; HAFNER & LYON, \textit{supra} note 2, at 61.
\textsuperscript{29} ABBATE, \textit{supra} note 6, at 15.
\textsuperscript{30} \textit{Id.} at 17.
\textsuperscript{31} \textit{Id.}
Although Baran’s proposal yielded substantial economic benefits, enhancing survivability represented the heart of his approach.\(^\text{33}\) Indeed, other aspects of Baran’s design underscore his preference for survivability over economic benefits. For example, Baran proposed that the nodes be located far away from population centers and be built with substantial excess capacity.\(^\text{34}\) He also advocated that the network incorporate both cryptography to ensure security, and a priority system to allow messages from high-level users to take precedence over those from lower-level users.\(^\text{35}\)

B. The Partial Reception of Baran’s Vision

Despite Baran’s best efforts, his network was never built. Although in August 1965 he convinced the Air Force to construct a system along the lines he proposed, the Department of Defense (DOD) reassigned responsibility for building and running the network to the Defense Communications Agency, which lacked any expertise in digital technology.\(^\text{36}\) Rather than see such inexperienced hands mismanage the network, Baran and his Air Force supporters opted to kill the project, and Baran moved on to other matters. The eventual construction of a network based Baran’s principles would have to wait until 1969, when the DOD created the ARPANET.\(^\text{37}\)

Meanwhile, across the Atlantic, UK researcher Donald Davies independently developed the idea of packet switching shortly after Baran. Although Davies retrospectively acknowledged

\begin{itemize}
  \item \(^\text{33}\) ABBATE, supra note 6, at 18.
  \item \(^\text{34}\) Id. at 11.
  \item \(^\text{35}\) Id.
  \item \(^\text{36}\) Id. at 20.
  \item \(^\text{37}\) Id. at 20–21; HAFNER & LYON, supra note 2, at 64.
\end{itemize}
Baran’s contributions, his primary motivation differed from Baran’s. Indeed, Davies sought to improve the economic efficiency of the network, not enhance its survivability in the face of a nuclear attack. Toward that end, he rejected the high levels of redundancy that Baran envisioned in favor of a more modest approach that integrated, among other things, a time-sharing principle.

Although Davies’s attempt to construct his network also failed for lack of funding, he was able to introduce both his and Baran’s work to Lawrence Roberts, who was in charge of building the ARPANET and had the budget to construct a network that incorporated new principles. However, Roberts’s goals, like Davies’s, differed from Baran’s, which caused him to adopt only part of Baran’s design. Most importantly, Roberts, like Davies, did not share Baran’s focus on survivability, as that was much less of a concern in nonmilitary contexts, such as research and commercial deployment. Instead, Roberts’s primary goal was to interconnect the several major research computing centers and the time-sharing computers that the Advanced Research Projects Agency (ARPA) was funding. Thus, Baran’s ideas about survivability “were

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38 ABBATE, supra note 6, at 27; HAFNER & LYON, supra note 2, at 65, 77. Interestingly, it was Davies who coined the terms packets and packet switching. Baran instead referred to message blocks and distributed adaptive message block switching. Davies would later say to Baran, “Well, you may have got there first, but I got the name.” HAFNER & LYON, supra note 2, at 67.
39 HAFNER & LYON, supra note 2, at 66.
40 ABBATE, supra note 6, at 27–29.
41 Id. at 29–33.
42 Id. at 37, 39, 41.
43 HAFNER & LYON, supra note 2, at 76.
44 Id. at 77.
all but ignored by the military” and “the topology of today’s Internet has little to do with his vision.”

Although Roberts rejected the high levels of redundancy, reliability, and excess capacity that Baran’s distributed network required, he did adopt Baran’s emphasis on decentralized autonomous routing and packet switching. Indeed, the primary goals underlying the ARPANET’s design were to keep average delay time below 0.2 seconds and decrease cost; reliability, along with flexibility and expandability, were secondary concerns. In addition, under the ARPANET’s design, reliability was defined as requiring at just two nodes to fail before the network would become disconnected—a stark contrast to Baran’s focus on designing a network that could continue to operate even if as many as half of its nodes were destroyed.

Figure 2 depicts the manifestation Roberts’s work: a historical representation of the ARPANET design from 1969 to 1977, when it had reached maturity. A casual look at Figure 2 shows that the vast majority of nodes were connected by only two links, and a more in depth historical analysis confirms that between 1969 and 1977, the average number of links connected to the ARPANET nodes ranged from only 2.00 to 2.46, a far cry from Baran’s vision of highly redundant connections that would permit the network to operate effectively, even if it lost as many as half of its nodes.

46 BARABÁSI, supra note 2, at 144; accord CHRISTOS J.P. MOSCHOVITIS ET AL., HISTORY OF THE INTERNET: A CHRONOLOGY, 1843 TO THE PRESENT 35 (1999) (“Although Baran’s work was extremely influential on the ARPAnet founders, his imagined network never came to fruition.”).
47 ABBATE, supra note 6, at 39, 226 n.38.
49 Id.; see also ABBATE, supra note 6, at 57–58.
50 See Figure 2
51 See Figure 2
53 Baran, supra note 1, at 3.
In 1986, the National Science Foundation (NSF) began to deploy the NSFNET in order to provide remote access to supercomputing centers.54 This time, however, the federal government funded the project to fend off the danger of losing out to Japan in supercomputing leadership.55 By the end of 1989, the NSFNET had replaced the ARPANET.56

Source: Heart et al., supra note 52, at III-79 fig.5, III-81 fig.7, III-84 fig.10, III-89 fig.15.

54 Rogers, supra note 21, at 217.
55 Id.
56 HAFNER & LYON, supra note 2, at 255–56.
Unlike the ARPANET, university consortia or state-university partnerships often helped operate the NSFNET’s network through regional connections, which provided connectivity from distant locations to the NSFNET backbone’s interconnection points. In addition to being linked together, each of the supercomputing centers became the hub for a web connecting local campus networks to other universities. Notwithstanding the operational differences, however, the topology the NSFNET’s network ultimately looked similar to that of the ARPANET’s. The average number of links connected to each supercomputer center was only 2.38. Furthermore, the NSFNET backbone, portrayed in Figure 3, is functionally equivalent to the network of nodes connected by thicker white lines that appear at the top of Figure 4; the main difference is that

58 See Figure 4
59 See Figure 3.
60 Calculated from Figure 3.
they are placed with only a slight slant backwards. The links to local campuses are depicted by the thinner white lines emanating from each of those nodes.

**FIGURE 4: THE TOPOLOGY OF THE NSFNET’S REGIONAL NETWORKS**


The NSFNET’s topology required every packet traveling through the network to traverse a series of three-level hierarchies. 61. Specifically, a campus network would hand off the traffic it was originating to its regional network. 62 The regional network would then hand the traffic off to the NSFNET backbone. 63 The backbone would then route the traffic to the appropriate

61 See Figure 5
62 YOO, supra note 57, at 57.
63 Id.
regional network.\(^6^4\) That regional network would then pass the packets along to the campus network that was the ultimate destination.\(^6^5\) This routing architecture permitted each campus network to connect to only a single regional network.\(^6^6\)

**Figure 5: The NSFNET Three-Tiered Network Architecture**

![Diagram of the NSFNET three-tiered network architecture](source:image)

Source: Yoo, *supra* note 11, at 83 fig.2.

The federal government retained this three-tiered hierarchy when it privatized the NSFNET in 1995.\(^6^7\) The result: the NSFNET backbone was replaced by a group of Tier-1 Internet Service Providers (ISPs)—also called backbones—interconnecting four network access points. The regional networks were replaced by Tier-2 ISPs—also called regional ISPs.\(^6^8\) The

\(^{6^4}\) *Id.*  
\(^{6^5}\) *Id.*  
\(^{6^6}\) *Id.*  
\(^{6^7}\) *Id.* at 58.  
\(^{6^8}\) *Id.* at 58–59.
campus networks were replaced by Tier-3 ISPs—also called last-mile providers. Figure 6 shows that the privatized network can more effectively be illustrated through concentric circles, where the core of the network is an amalgamation of interconnected backbones, rather than though the more vertical configuration that represented the NSFNET’s three-tiered network, as depicted in Figure 5.

**FIGURE 6: THE NETWORK HIERARCHY DEPICTED AS A SERIES OF CONCENTRIC RINGS**

![Network Diagram](image)

Source: Yoo, *supra* note Error! Bookmark not defined., at 59 fig.4-4.

Comparing the concentric network in Figure 6 to the three networks that Baran laid out—see Figure 1—demonstrates that the architecture of the NSFNET’s privatized network bears greater similarity to the decentralized network that Baran disfavored rather than the distributed

69 See Figure 6
network that he preferred.\textsuperscript{70} This three-tiered hierarchy has become the standard reference model taught in every computer networking textbook.\textsuperscript{71}

II. \textbf{NETWORK SCIENCE’S INSIGHTS INTO THE BENEFITS OF NETWORKS}

As a historical matter, neither the ARPANET nor the NSFNET embraced Baran’s vision of a distributed, lattice-like network.\textsuperscript{72} Consequently, neither did the Internet.\textsuperscript{73} Far from being a mere matter of path dependence, the burgeoning field of network science, pioneered by scholars such as Duncan Watts, Albert-László Barabási, and others, has instead provided a present-day justification for adopting the decentralized, hierarchical Internet structure.

\begin{itemize}
\item \textsuperscript{70} Compare Figure 6 and Figure 1(b)
\item \textsuperscript{71} Dhamdhere & Dovrolis, supra note 10, at 1.
\item \textsuperscript{72} See supra notes 39–40, 43–45 and accompanying text.
\item \textsuperscript{73} See supra notes 67–71 and accompanying text
\end{itemize}
A. Duncan Watts and Small-World Networks

One of the most enigmatic principles in network science is commonly known as the small-world phenomenon. Specifically, common sense would suggest that as the size of a network increases, the number of links that one would need to traverse in order to connect two arbitrary nodes should increase as well. However, the works of Stanley Milgram and others have suggested otherwise. In his famous study, Milgram tested how many links were necessary to connect a randomly selected person—or “node”—in Kansas or Nebraska to a specific person in the Boston area. Despite speculation that such a connection might require 100 or more links, Milgram found that the chains connecting these two people traversed an average of only 5.7 links, a finding that was later popularized by the play, *Six Degrees of Separation*, which explored the notion that any two people in the world can be connected through fewer than six other people. In addition, Milgram found that certain key intermediaries—most notably in the case of one study a clothing merchant, [example of intermediaries]—played an outsized role in establishing the connection between Nebraska and Boston.

Duncan Watts, initially with Steven Strogatz in their landmark article, *Nature*, and later in his book, *Six Degrees*, sought to further synthesize the principles underlying the small-world phenomenon by providing an explanation of how the random introduction of a small number of

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75 Milgram’s initial experiment measured how many links it took to connect volunteers in Wichita, Kansas, to the wife of a divinity school student living in Cambridge, Massachusetts. *Id.* at 64. Milgram’s second experiment, conducted with Jeffrey Travers, measured how many inks it too to connect volunteers in Omaha, Nebraska, to a stockbroker working in Boston and living in Sharon, Massachusetts. *Id.*; Jeffrey Travers & Stanley Milgram, *An Experimental Study of the Small World Problem*, 32 SOCIOMETRY 425, 429–30 (1969).
76 Milgram, supra note 74, at 65.
77 Travers & Milgram, supra note 75, at 436–38 & tbl.2.
79 Milgram, supra note 74, at 66; Travers & Milgram, supra note 75, at 436–39.
long distance connections can allow highly clustered networks to still retain short path lengths.\textsuperscript{80} Taking a cue from Paul Erdős and Alfred Rényi, Watts considered at one extreme those networks where the links between the nodes are completely random.\textsuperscript{81} Random networks are characterized by two notable qualities. First, because long distance and short distance links appear with the same probability, random networks exhibit very low levels of local clustering.\textsuperscript{82} Second, the presence of long-distance links means that the typical path will grow relatively slowly with the size of the network, with the average path length rising logarithmically with the number of nodes.\textsuperscript{83}

These random networks stand in stark contrast to both Baran’s distributed network and other lattice-like networks, in which every node is connected to each of its immediate neighbors. Indeed, highly ordered, lattice-like networks exhibit different qualities than random networks. First, because their nodes are connected exclusively by short distance links to their immediate neighbors, traditional lattice-like networks exhibit a high level of clustering.\textsuperscript{84} Second, the absence of long-distance links means that the average path length will grow linearly with the size of the network, which is much faster than logarithmic growth.\textsuperscript{85} Finally, the longer average path length associated makes lattice-like networks inherently \textit{large-world networks}; specifically, the number of links needed to travel between two randomly chosen nodes in the network tends to be rather large. The shorter path length associated with random networks, in contrast, makes them

\textsuperscript{82} Watts & Strogatz, \textit{supra} note 80, at 440.
\textsuperscript{83} Id.
\textsuperscript{84} Id.
\textsuperscript{85} Id.
inherently small-world networks, in that the number of links needed to travel between two randomly selected nodes tends to somewhat small.  

To demonstrate this difference, Watts and Strogatz compared the behavior of a lattice-like network with the that of a random network.  

To avoid complications created by nodes located on the boundary of the graph, which necessarily have fewer interactions than other nodes,, Watts and Strogatz modeled the lattice-like network as a circle, with each node connected to its two nearest neighbors and not to any distant nodes.  

To be sure, this conception of the lattice-like network differs from Baran’s distributed network, in which each node should be connected to four neighboring nodes.  

Watts and Strogatz then ran an experiment where each link could be reassigned with probability $p$ to another randomly chosen node in the circle.  

If the probability was zero, no links would ever be reassigned. Under that condition, the network should remain both a strict lattice and a large-world.  

If the probability was 100 percent, all links would be reassigned, and the network should become both an Erdös-type random network and a small-world network.

86 Id.
87 Id.
88 WATTS, supra note 80, at 84; see the left-hand graph in Figure 7
89 See supra note 25 and accompanying text.
90 Watts & Strogatz, supra note 80, at 440.
91 See the left-hand graph of Figure 5.
92 Watts & Strogatz, supra note 80, at 441.; See the right-hand graph of Figure 7—
A relevant takeaway from Watts and Strogatz’ work is what happens when only a small number of links is reassigned, as depicted in the middle graph of Figure 7. Specifically, Watts and Strogatz found that introducing a small number of long distance links into the network radically reduced the average path length, while still retaining a highly clustered structure.\(^{93}\)

This showed how short cuts, as Watts and Strogatz called them, can allow networks to be both highly clustered and small-world networks.\(^{94}\) Indeed, the introduction of five such short cuts reduced the average path length by half.\(^{95}\) The effect does decay rather rapidly, however. A further 50 percent reduction in path length would require introducing an additional 50 short cuts.\(^{96}\) Watts and Strogatz found short cuts to be empirically present in a wide range of

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\(^{93}\) Watts & Strogatz, supra note 80, at 440–41.

\(^{94}\) Id. at 441 fig.1 & 2.

\(^{95}\) Watts, supra note 80, at 89

\(^{96}\) Id.
networks, including those of film actors, the electrical power grid, and the neural network of a particular nematode worm.97

The implications of these findings are profound. Watts and Strogatz showed that adding a limited number of random shortcuts to a strict lattice-like structure can play a key role in permitting Baran’s distributed networks to scale.98 Under this model, each node is no longer undifferentiated in terms of connectivity. Instead, there are a handful of nodes with more important levels of connectivity that keep the network small-world, which enables the network to scale, a characteristic that had eluded Baran’s distributed network. A notable caveat, however, is that these nodes start to increasingly resemble the centers of the smaller stars in Baran’s decentralized network.

B. Albert-László Barabási and Scale-Free Networks

Watts’s work was extended by Albert-László Barabási, who noted that all of the nodes in the random networks in the Erdős-Rényi and Watts-Strogatz studies had a functionally equivalent number of links.99 In other words, the number of links connected to the typical node in a random network has a given scale, a specific characteristic that is represented across all nodes in the network.100 In random networks, the number of links associated with each node should thus follow a Poisson distribution, which should occur when all nodes have roughly the

97 Watts & Strogatz, supra note 80, at 441 & tbl.1.
98 See supra note 93–94 and accompanying text.
99 Reka Albert & Albert-László Barabási, Diameter of the World-Wide Web, 401 NATURE 130, 130 (1999);
100 Barabási et al., supra note 99, at 178.
same number of links. However, Barabási observed that the Internet evinced a very different pattern. Instead of following a Poisson distribution, the empirical literature indicated the number of links connected to each node exhibited a power-law distribution, which occurs when there are a small number of nodes with a disproportionately large number of links. The heterogeneity in the number of links attached to each node in such networks meant that the network did not have an inherent characteristic of scale in terms of number of links. As such, Barabási called such networks *scale-free* networks.

Barabási proceeded to develop a dynamic theory to explain why scale-free networks would develop under the Watts-Strogatz and other similar models. Specifically, Barabási observed that those models assumed that the number of nodes in a given network is static, and network connections form randomly among nodes. With that observation, Barabási created a model with a dynamic network, where new nodes would be free to decide where to attach to the network. Under Barabási’s model, these nodes would follow the practice of *preferential attachment*: they would be more likely to attach to nodes that were already highly connected.

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104 *Id.* at 510–11.
105 *Id.* at 511.
106 *Id.*
Barabási’s result: a network where the majority of nodes have only one or two links while a handful of other nodes—those that attracted the new nodes—have a large number of links.\textsuperscript{107}

The presence of these \textit{hubs} radically reduces the shortest path between two nodes.\textsuperscript{108} Moreover, in a random network the average path length should grow linearly with the number of nodes in the network, but Barabási’s work suggested that the average path length would grow logarithmically—and thus considerably slower—in scale-free networks.\textsuperscript{109} He noted that a wide range of networks had empirically followed this pattern, including the Internet, the World Wide Web, film actor collaborations, scientific collaborations, human sexual contacts, cellular networks, ecological networks, phone call networks, citation networks, linguistic networks, power networks, neural networks, and protein folding.\textsuperscript{110} Later empirical studies refined Barabási’s, particularly as it applied to the Internet. Specifically, these studies found that the Internet’s architecture did not resemble a lattice-like grid and that the preferential attachment rates fell slightly below that identified in Barabási’s hypothesis.\textsuperscript{111} Other studies also found preferential attachment, but concluded that preferential attachment occurred at rates above those Barabási predicted.\textsuperscript{112}

Germane to Baran’s work, Barabási pointed out that decentralized networks with a limited number of hubs are actually \textit{more} resilient against incidental node and link failures; indeed, random failure are almost certainly likely to affect small nodes, which have little to no

\begin{itemize}
\item \textsuperscript{107} \textit{Id.}
\item \textsuperscript{108} Albert-László Barabási, \textit{The Physics of the Web}, PHYS. WORLD, July 2001, at 33, 34 & fig.1.
\item \textsuperscript{109} Albert et al., \textit{supra} note 99, at 130.
\item \textsuperscript{110} Albert & Barabási, \textit{supra} note 102, at 49–54.
\item \textsuperscript{112} Qiang Chen et al., \textit{The Origin of Power Laws in Internet Topologies Revisited}, 2 PROC. IEEE INFOCOM 608 (2002).
\end{itemize}
impact on the network’s overall connectivity. He did note, however, that the presence of key focal points makes decentralized networks more vulnerable to *deliberate attacks*, which would undoubtedly target the hubs.

C. Implications for the Distributed Network

The developments in network theory provide insight into why the Internet did not incorporate the topology that Baran envisioned. The fact that average path length in traditional lattice-like networks grows linearly with the number of nodes means that such networks are difficult to scale. The introduction of long-distance short cuts and hubs allows average path length to grow at a slower, logarithmic rate, as depicted in Figure 8.

**Figure 8: Average Path Length Growth for Distributed and Decentralized Networks**

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115 *See supra* note 85 and accompanying text.
Hubs thus enable shortcuts to play a critical role in giving networks the small-world characteristics in which average path length remains intact as the network grows. These features deviate from Baran’s distributed structure—in which all nodes are undifferentiated and of equal importance—and instead make certain nodes more important. Network theory thus provides an explanation of why the Internet mimicked a more decentralized structure of satellite hubs instead of Baran’s distributed structure. Some commentators have raised the concern that the increased importance of some hubs may give them a competitive advantage.116 But such concerns can be alleviated by adding competition from new hubs so long as the costs of network reconfiguration are manageable.

III. THE EVOLUTION OF THE MODERN INTERNET

The conventional wisdom that the Internet’s foundational topology adhered to a three-tier hierarchy was so well established that it served as the reference model in every textbook on computer networking. As noted above, the resulting architecture bore a far stronger resemblance to what Baran called a decentralized network than it did to his distributed network.

More recently, two forces have pushed the network even further away from Baran’s distributed model. The first is the triumph of the client-server architecture over the peer-to-peer model, as discussed in Part A. The second is the emergence of topologies that deviate from the three-tiered hierarchical model, as discussed in Part B. Although changes such as secondary peering and multihoming may make the network more mesh-like, other features make the hierarchy even flatter (such as content delivery networks and direct interconnection between content providers and last-mile networks). Interestingly, these changes have made the architecture begin to resemble the star-like network that Baran described as centralized.

A. The Rise of the Client-Server Architecture and the Decline of Peer-to-Peer

During the Internet’s early days, before it became a mass market phenomenon in the mid-1990s, it was used primarily by tech-savvy researchers to send email, share files, and view bulletin boards. When that was the case, the network could be characterized as a “peer-to-peer any-to-any flat mesh” in which nodes would typically request and send data. Using the early

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117 See supra note 71 and accompanying text.
118 Christopher S. Yoo, Free Speech and the Myth of the Internet as an Unintermediated Experience, 78 GEO. WASH. L. REV. 697, 703 (2010).
network in that way was consistent with the assumption of Baran’s distributed network that
nodes were generally homogeneous and of equal status.

The advent of the World Wide Web during the mid-1990s upended this dynamic.\textsuperscript{120} The
Internet began to embrace a client-server architecture, which was inherently asymmetric: nodes
were divided into the end users who requested content and the content providers who furnished
it.\textsuperscript{121} Although file sharing slowed this transformation for a time, peer-to-peer traffic has been
more or less steadily declining as a percentage of traffic since at least 2005,\textsuperscript{122} and has been
decreasing in absolute terms since 2012.\textsuperscript{123} Cisco forecasts that it will continue to decline, sinking
to a mere 2\% by 2021.\textsuperscript{124}

\textbf{FIGURE 9: PEER-TO-PEER TRAFFIC AS A PERCENTAGE OF CONSUMER INTERNET TRAFFIC}

\begin{figure}
\centering
\includegraphics[width=\textwidth]{peer-to-peer_traffic_percentage.png}
\end{figure}

\begin{enumerate}
\item Yoo, \textit{supra} note 118, at 703.
\item \textit{Id.} Huston, \textit{supra} note 119, at 4.
\item See Figure 9; Yoo, \textit{supra} note 57, at 30–31.
\end{enumerate}

The asymmetry inherent in the client-server architecture is inconsistent with Baran’s vision of a distributed model with homogeneous nodes. The effect of this asymmetric architecture has been compounded as traffic has become increasingly concentrated among a small number of content providers. In such an environment, it no longer makes business sense for content providers to reside at the edge of the network and rely on a series of public networks to transmit their data to end users, particularly when there are delays associated with DNS resolution, TCP and TLS handshakes, and long round-trip times. In addition, content providers can avoid the uncertainties inherent in sharing bandwidth with others, and the concomitant possibilities of delay and security risks, by building their own private networks out to local caches. When that is the case, it makes sense for content providers to bypass the network and host their content as close to the end users as possible.

This dynamic explains why large content providers have begun to connect directly to last-mile providers. Indeed, in 2008 Google had already begun to construct a wide area network that covered most of the U.S. and connected to Europe, Asia, and South America. Yahoo! and
Microsoft were likewise constructing wide area works covering the U.S. More recent press reports have covered how that Google, Facebook, Microsoft, and Amazon are currently constructing large networks of undersea cables to interconnect proprietary data centers located ever closer to end users. In fact, the international capacity of these companies increased tenfold from 2013 to 2017. The upshot: rather than waiting for users to reach out for content, content providers are increasingly bringing content directly to end users.

B. The Flattening of the Network Architecture

The architecture has evolved considerably over the past three decades. In the words of the emerging empirical literature, the Internet’s architecture has become flatter; traffic is no longer passing through all three layers of the Internet’s traditional three-layer hierarchy.

This flattening of the architecture may be the result of different causes. First, it may be from increased density of connections in the core, either from Tier 3 ISPs bypassing Tier 2 ISPs and connecting directly to Tier 1 ISPs or from Tier 2 ISPs bypassing Tier 1 ISPs by engaging in secondary peering. If this is the case, the network can be said to be becoming more mesh like. Alternatively, large content providers could be bypassing both Tier 1 and Tier 2 ISPs altogether motivated by a 2005 press report about the growth of Google’s private network. R. Scott Raynovich, Google’s Own Private Internet, LIGHT READING (Sept. 20, 2005), https://www.lightreading.com/ethernet-ip/googles-own-private-internet/d/d-id/618122.


Id. See Christopher S. Yoo, Innovations in the Internet’s Architecture that Challenge the Status Quo, 8 J. ON TELECOMM. & HIGH TECH. L. 79, 85–90 (2010) (discussing how that architecture of the internet has evolved). See Gill et al., supra note 131; Besen & Israel, supra note 10; Dhamdhere & Dovrolis, supra note 10. YOO, supra note 58, at 64–68.
and using content distribution networks and data centers to connect directly to Tier 3 ISPs. If this is the case, the network is more properly described as becoming more centralized.

A key indicator of a network’s true underlying architecture is the rate of change in the average path length over time. Specifically, and as noted earlier, if the network is distributed, average path length should grow linearly with the number of nodes. If the network is decentralized, average path length should grow logarithmically with the number of nodes. `associated with centralized networks. In the limit, the average path length of a centralized star network is two.

Empirical studies of changes in the average path length thus provide important insights into how the architecture is changing. If the network is decentralized or distributed, the fact that the number of nodes in the network has been constantly increasing implies that the average path length should be growing. However, and somewhat surprisingly, empirical studies have found the contrary: the average path length in the Internet has been either flat or decreasing, despite the growth in the number of nodes. This phenomenon suggests that the network is not becoming more mesh like; it is instead becoming more centralized. Under this scenario, the flattening of the Internet’s architecture is more associated with content providers connecting directly with Tier 3 ISPs on the edge of the network rather than the result of denser peering.

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137 See supra notes 80, 85–90, 93–95, 109 and accompanying text.
138 See supra note 86 and accompanying text. See supra note 109 and accompanying text.
139 See supra note 109 and accompanying text.
relationships in the core of the network. In other words, the flattening of the architecture begins to more closely reflect Baran’s centralized network, where the average path length is 1.

Studies have attempted to move beyond merely studying average path length by examining directly whether the flattening of the network is associated more with the densification of the core or with the ISP’s bypassing the core altogether. The earliest studies observed that traffic from the largest content providers traversed fewer hops on Tier 1 networks than did traffic from smaller content providers.142 Although a 2011 study found it plausible that an increase in multihoming, rather than content providers bypassing Tier 1 ISPs,143 was the cause of the lack of growth in average path length, a different 2011 study—one that actually evaluated the network empirically—found otherwise; the latter reinforced the conclusion in the earlier studies: the flattening of the network has stemmed from content provider forging routes that bypass Tier 1 ISPs rather than form multihoming.144 What is more, a 2012 study implicitly minimized the impact of increased connectivity among Tier 1 and Tier 2 ISPs when it concluded that although the core of the network was densifying, the interconnectedness of the overall network remained unchanged.145

Together these studies contradict suggestions that the flattening of the Internet means that the network is becoming more mesh like, and instead indicate that content providers are increasingly bypassing the other tiers and connecting directly to last-mile ISPs. This in turn suggests the possibility that the standard reference model should be replaced with a radically different vision: a stack of star networks as depicted in Figure 10, each centered on one of the

142 Gill et al., supra note 131, at 6; Craig Labovitz et al., Internet Inter-Domain Traffic, Proc. ACM SIGCOMM 2010 Conf. 75, 79 (2010).
143 Dhamdhere & Dovrolis, supra note Error! Bookmark not defined., at 1424–25.
144 Xiang et al., supra note Error! Bookmark not defined., at 3–4.
145 Li et al., supra note Error! Bookmark not defined., at 5.
leading content providers and connecting by private links (represented by dotted lines) to the same public locations at the edge (represented by solid lines). Rather than resembling a mesh network, this emerging architecture bears greater similarity to Baran’s centralized network.

**FIGURE 10: NETWORK ARCHITECTURE RECONCEPTUALIZED AS A STACK OF STARS CENTERED ON CONTENT PROVIDERS**

This reconceptualization has profound implications. Investment in network capacity helps meet users’ ever-growing demand for bandwidth. Moreover, content providers and their customers can benefit from improved risk management and cost control, support for the deployment of new applications, and improvements in quality of service.\(^{146}\) At the same time, as Geoff Huston has noted, the resulting topology would resemble the historical broadcast television network, in which local distribution networks that carried content from multiple

\(^{146}\) Yoo, *supra* note 57, at 68; Gill et al., *supra* note 131, at 2–3, 7–8.
sources were fed by proprietary satellite networks. The resulting diversion of long-haul traffic from public to private networks may make backbones more difficult to sustain, which in turn may adversely affect smaller content providers that are unable to pursue the same strategy. Increasing centralization may increase the ability for some actors to serve as gatekeepers.

**CONCLUSION**

The seminal work of Paul Baran has never received its full due. Although his breakthrough research established the foundation upon which modern packet switching was built, Baran’s three-part taxonomy of centralized, decentralized, and distributed networks has been overlooked—the reception of his ideas remains incomplete. Specifically, the early Internet architects were less concerned with survivability than was Baran while network theorists such as Watts and Barabási showed hubs, shortcuts, and other modifications to the distributed architecture can allow networks to scale. The result was a three-tiered hierarchy of ISPs that formed what Baran would define as a decentralized network, not the distributed network that he preferred.

The dramatic changes to the network hierarchy over recent years suggests that we should reconceptualize the architecture still further. Specifically, the fact that the average path length is shrinking while routes are increasingly bypassing Tier 1 ISPs to connect directly to the last mile compels us to consider whether the standard reference model should be replaced with a radically

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different vision: a stack of star networks, each centered on one of the leading content providers, which would functionally bear a stronger resemblance to Baran’s centralized network.

The expansion of content providers into wide area networking may yield a number of substantial benefits. Greater investment in transmission capacity is essential to meet users’ ever-growing demand for bandwidth. However, the shift from public to private backbones may certain implications that need to be thought out. Such concerns must be analyzed in light of the full range of options instead of being simply asserted. The full impact of the change can only be understood once the full choice set is endogenized, including multihoming, the creation of alternative hubs, cooperative investments in competing technology, and other institutional forms, to name a just few.